

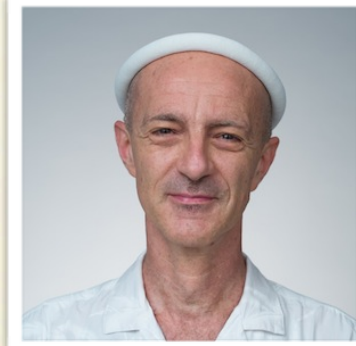
**INTERACTING HOPF ALGEBRAS**  
**THE THEORY OF LINEAR SYSTEMS**

**FILIPPO BONCHI**  
**UNIVERSITY OF PISA**

# Collaborators



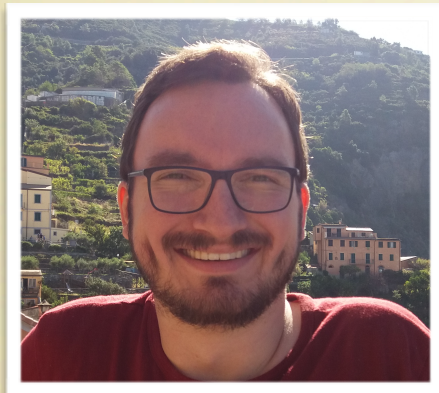
Fabio  
Zanasi  
UCL



Dusko Pavlovic  
Hawai'i



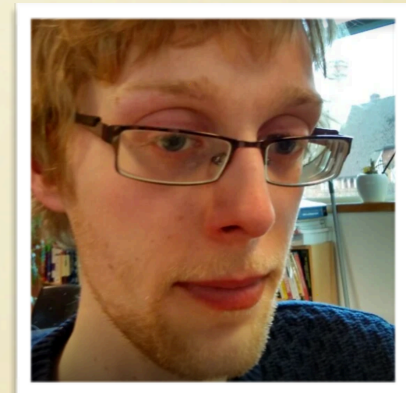
Pawel  
Sobocinski  
Southampton



Jens Seeber  
IMT Lucca



Robin Piedeleu  
Oxford



Josh Holland  
Southampton

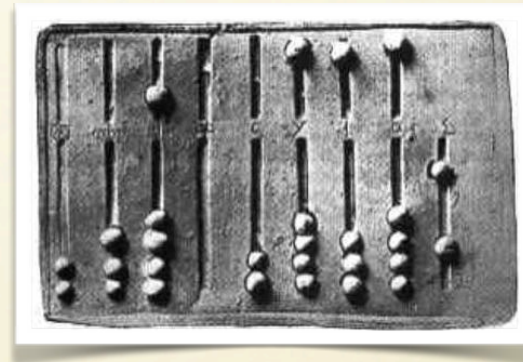
**I. INTERACTING HOPF ALGEBRAS**

**II. CONTROVERSIAL INTERMEZZO**

**III. RAMIFICATIONS**

# Numeral Systems

Roman numbers are quite inconvenient to perform the elementary arithmetic operations



In *Liber Abbaci* (1201), Fibonacci introduce in Europe the Arab numeral system together with the basic algorithm

In 1280, the city of Florence forbade the use of Arab *ciphers*

Nowadays, the introduction of Arab numbers is considered a fundamental moment in the History of Science

# Equations

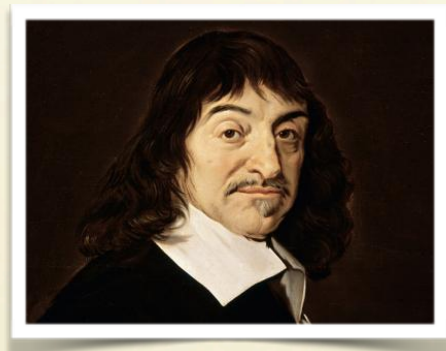
In *The Whetstone of Witte* (1557), Robert Recorde introduced the symbol =

$$14x + 15 = 71$$

Sign	Wording	Meaning
φ	dragma, numerus	the number, the absolute term
ⓧ	radix	the root, the unknown, x
z	zensus (census)	the square, $x^2$
∞	cubus	the cube, $x^3$
zz	zensdezens	the fourth power, $x^4$
β	sursolidum	the fifth power, $x^5$
z∞	zensicubus	the sixth power, $x^6$
bβ	bissursolidum	the seventh power, $x^7$
zzz	zenszensdezens	the eighth power, $x^8$
∞∞	cubus de cubo	the ninth power, $x^9$

# Polynomials

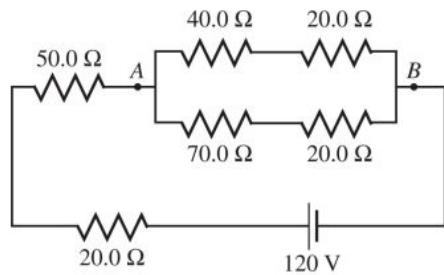
Our notation for polynomials was introduced by Descartes in *La Géométrie* (1637)



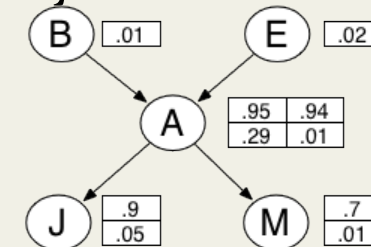
After we started writing  $x^2$ ,  $x^3$ , ...  $x^n$ ,  
we could think about  $x^{-2}$  or  $x^{0.5}$

# Network diagrams

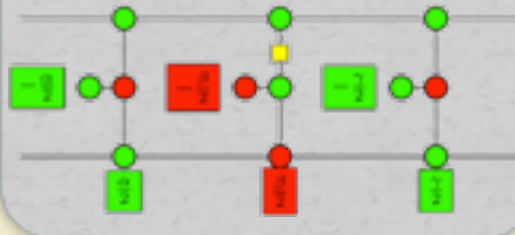
## Electrical Circuits



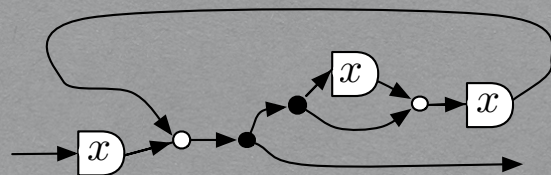
## Bayesian Networks



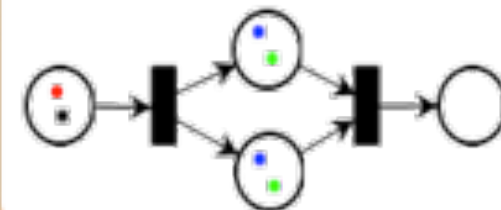
## Quantum Processes



## Signal Flow Graphs

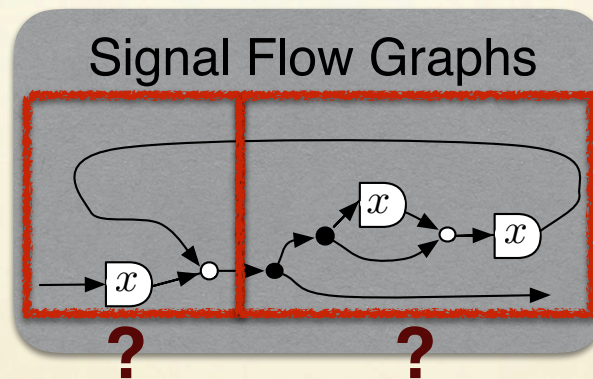


## Petri Nets



# Compositional Semantics

Diagrammatic languages are not really made of syntax.



We are able to describe the behaviour of the whole systems

But not the behaviour of the single components

The behaviour of the whole system should be "reducible" to the behaviour of its components



# Compositional Modelling

There is an emerging, multi-disciplinary field aiming at studying different sorts of networks **compositionally**, inspired by the **algebraic methods** of programming language semantics.



Diagrams are first-class citizens of the theory. The appropriate algebraic setting is **monoidal** (and not **cartesian**) categories.

# More and more influential

## The Future Will Be Formulated Using Category Theory



**Jayshree Pandya** Contributor  
**COGNITIVE WORLD** Contributor Group ⓘ

AI & Big Data

*Jayshree Pandya is Founder of Risk Group & Host of Risk Roundup.*

---

*A new approach to defining and designing systems is coming.*

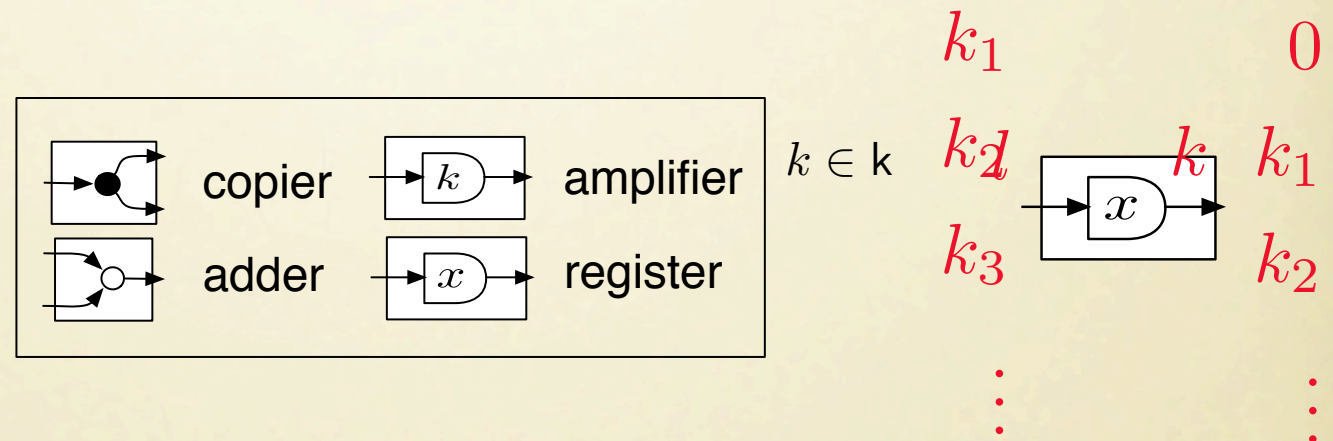
<https://www.forbes.com/sites/cognitiveworld/2019/07/29/the-future-will-be-formulated-using-category-theory/#71a09469625e>



# **I. INTERACTING HOPF ALGEBRAS**

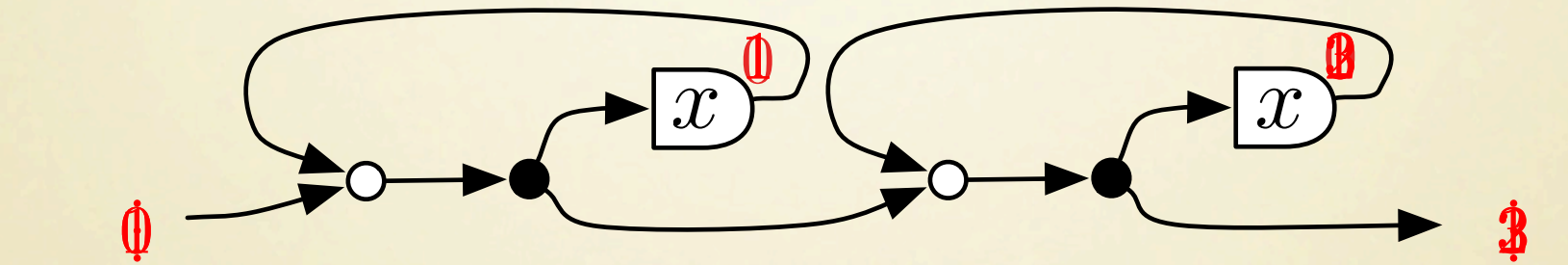
# Signal Flow Graphs

Signal Flow Graphs are **stream** processing circuits widely adopted in Control Theory and Signal Processing

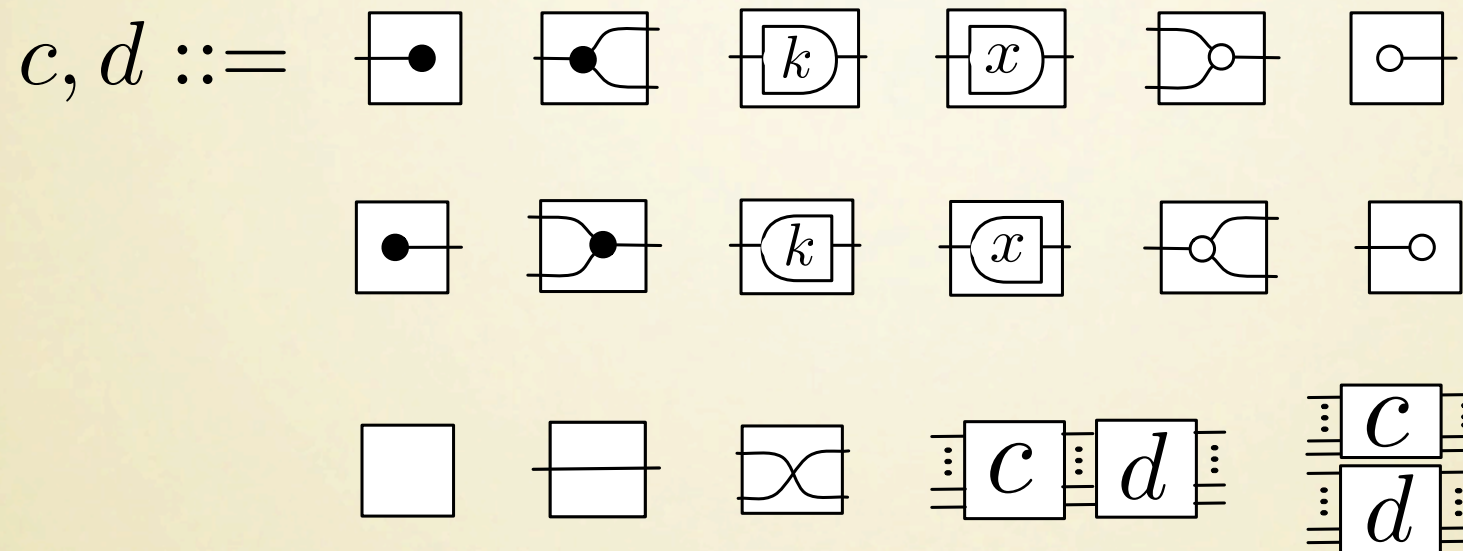


Claude Shannon. *The theory and design of linear differential equation machines* (1942).

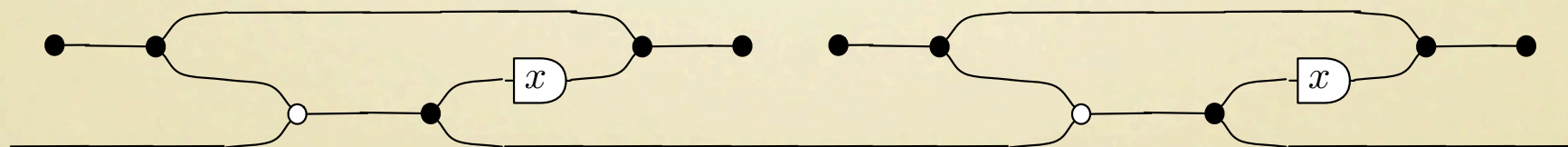
# Signal Flow Graphs



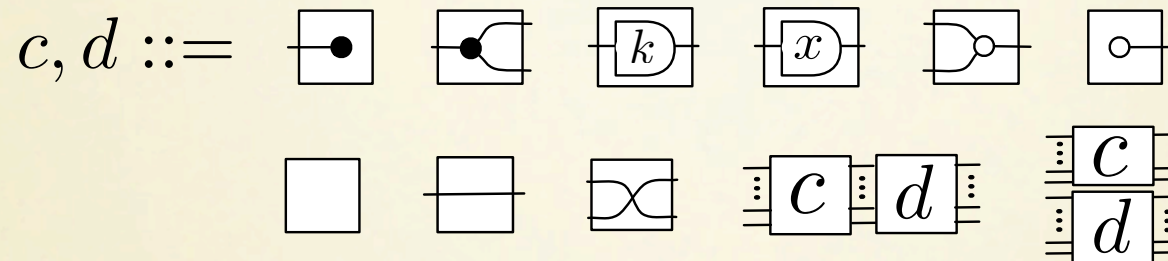
# String Diagrammatic Syntax



Subject to the laws of PROPs



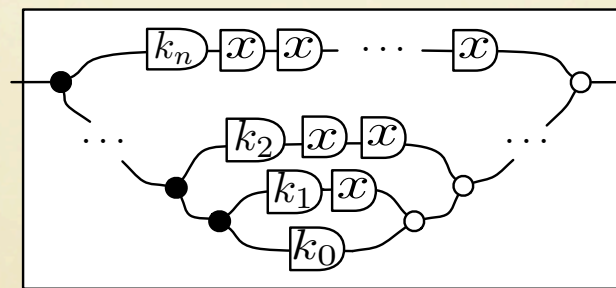
# Functional Circuits

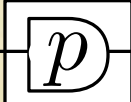


We can represent polynomial

$$p = k_0 + k_1x + \dots + k_nx^n$$

as



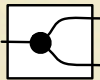
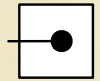
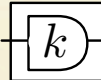

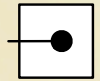
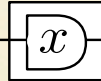




Hereafter denoted by 

# Polynomial Matrices

The denotational semantics of functional circuits is given by the PROP morphism

$$\llbracket \cdot \rrbracket : \overrightarrow{\text{Circ}} \rightarrow \text{Mat}_{k[x]}$$

defined inductively as follows

	$\mapsto$	$\begin{pmatrix} 1 \\ 1 \end{pmatrix}$		$\mapsto$	$!$		$\mapsto$	$(k)$		$\mapsto$	$\begin{pmatrix} 1 & 1 \end{pmatrix}$
	$\mapsto$	$!$		$\mapsto$	$(x)$		$\mapsto$	$?$			
	$\mapsto$	$id_0$		$\mapsto$	$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$		$\mapsto$	$id_1$			

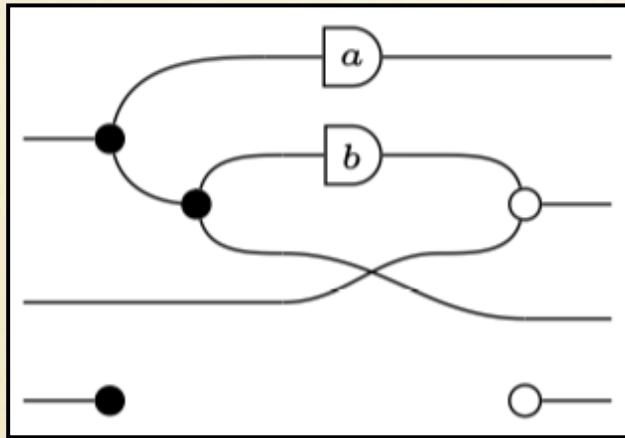
$$c_1 \oplus c_2 \mapsto \llbracket c_1 \rrbracket \oplus \llbracket c_2 \rrbracket$$

$$c_1 ; c_2 \mapsto \llbracket c_1 \rrbracket ; \llbracket c_2 \rrbracket$$

where  $!$  and  $?$  are the unique morphism given by initiality and finality of  $0$



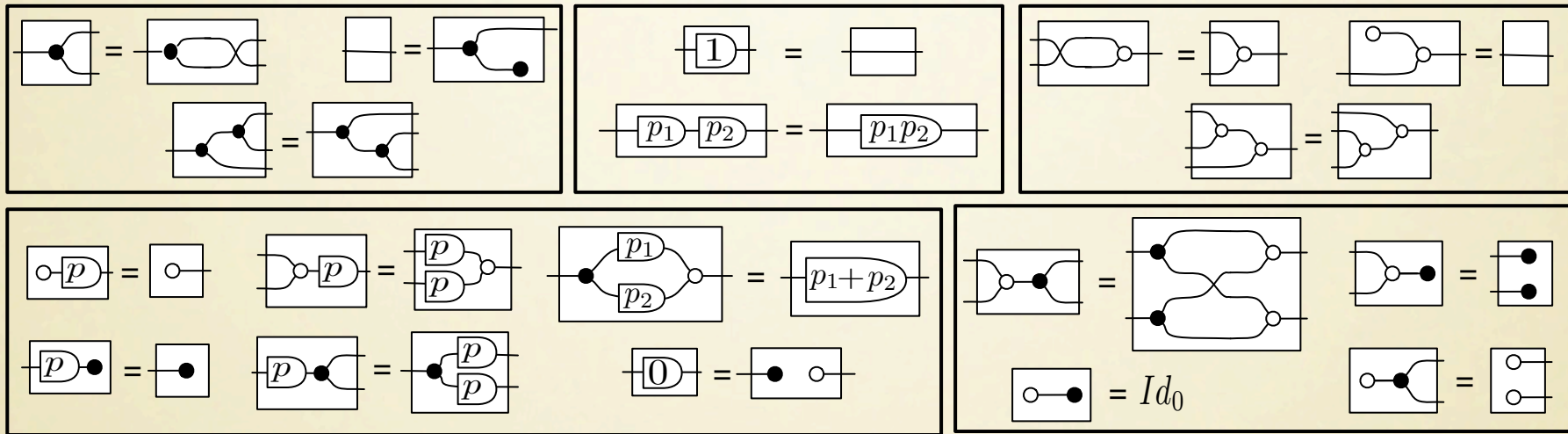
# Example



$$\begin{pmatrix} a & 0 & 0 \\ b & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

# Hopf Algebras

The theory  $\mathbb{H}\mathbb{A}$  is  $\overrightarrow{\text{Circ}}$  quotiented by the following axioms

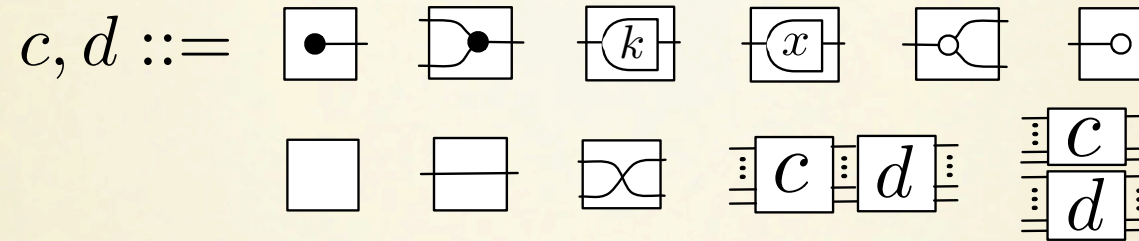


## Soundness and Completeness

$\mathbb{H}\mathbb{A}$  is isomorphic to  $\text{Mat}_{k[x]}$

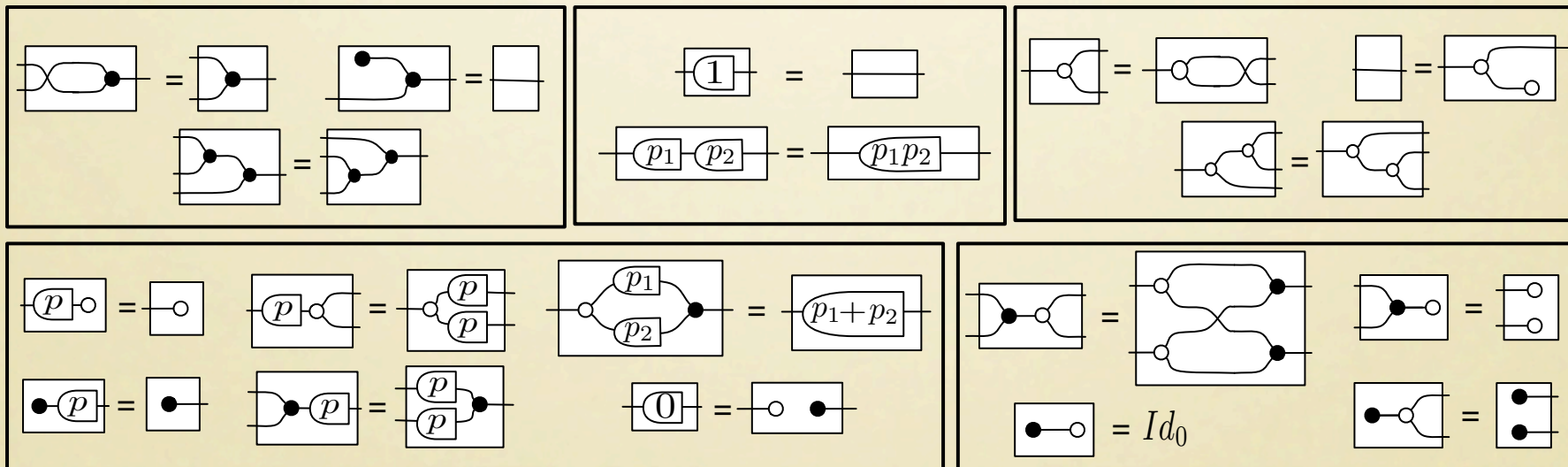
$$\overrightarrow{[c]} = \overrightarrow{[d]} \iff c \stackrel{\mathbb{H}\mathbb{A}}{=} d$$

# Cofunctional Circuits



The denotational semantics  $\llbracket \cdot \rrbracket : \overleftarrow{\text{Circ}} \rightarrow \text{Mat}_{k[x]}^{op}$  is given by duality

The theory  $\mathbb{H}\mathbb{A}^{op}$  is  $\overleftarrow{\text{Circ}}$  quotiented by the following axioms

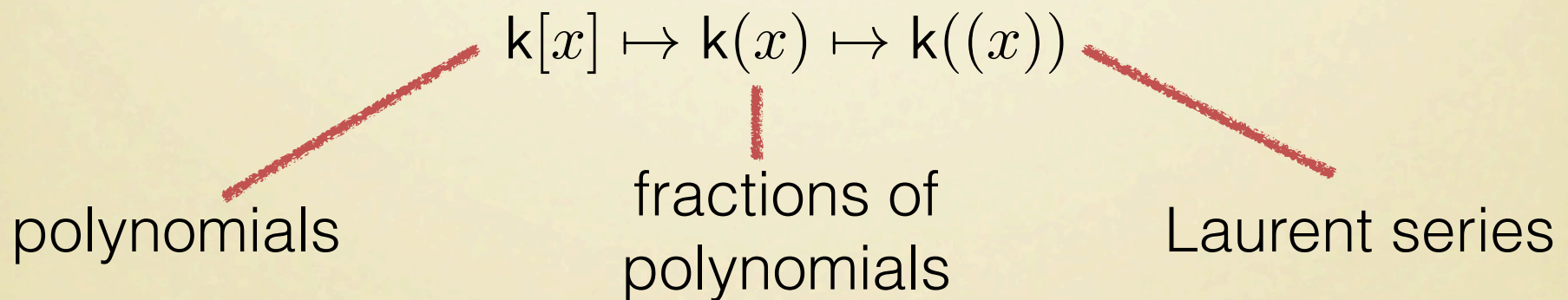


# Semantics of Generalised Circuits

When we allow combinations of functional and co-functional circuits (like in feedbacks) we may get *relational* behaviours:

For instance,  $\boxed{\bullet} \dashv \sigma ; \sigma \dashv \boxed{\bullet} \dashv \sigma$  expresses the diagonal relation

Moreover, polynomials are not enough:  
we need *fractions* of polynomials.

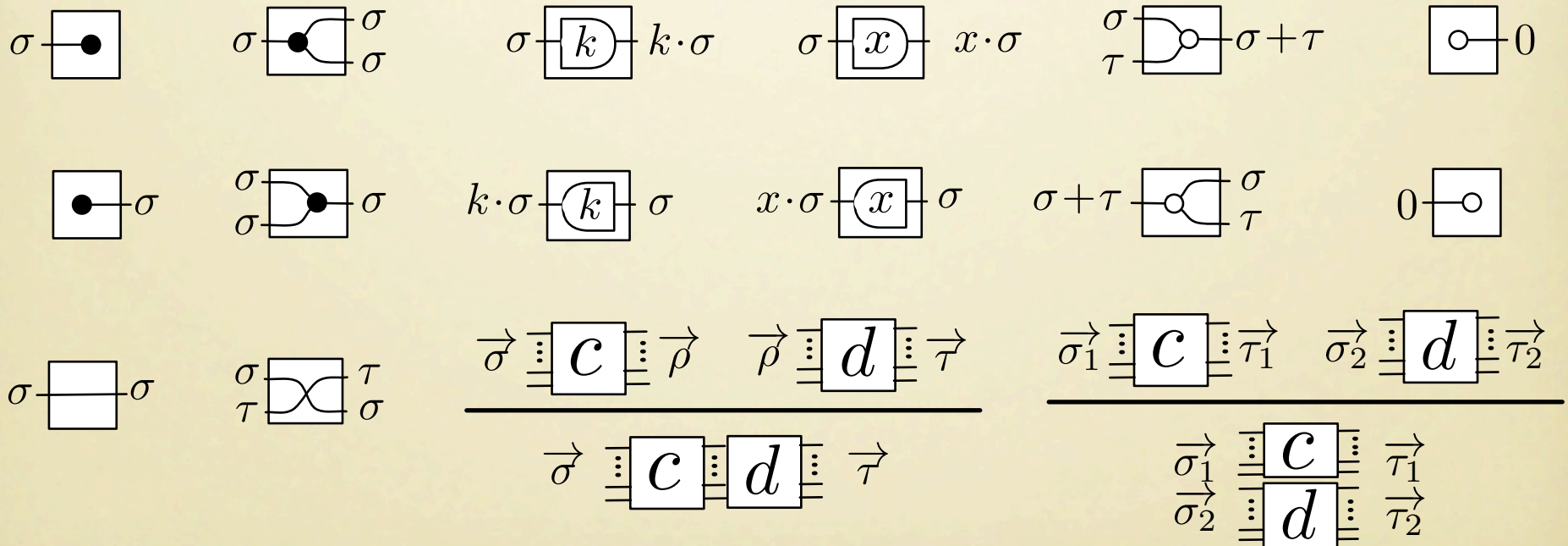


# Semantics of Generalised Circuits

The denotational semantics of circuits is given by the PROP morphism

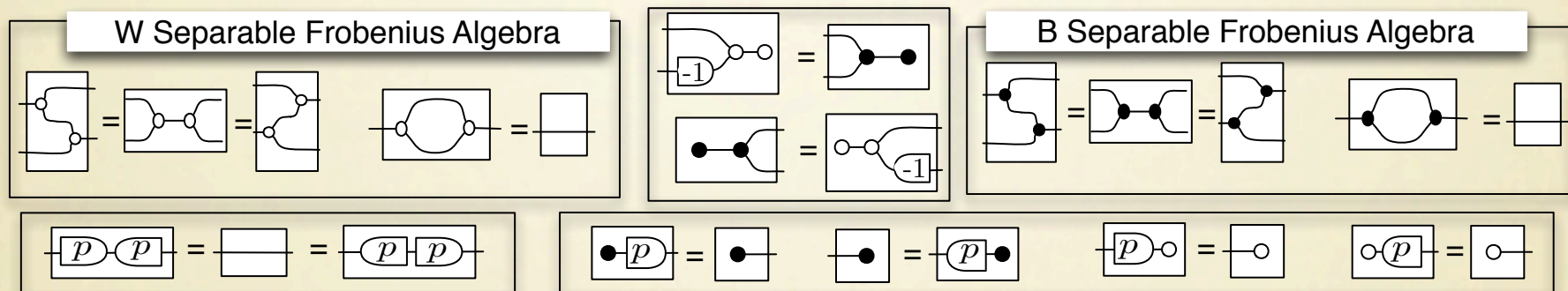
$$[[\cdot]] : \text{Circ} \rightarrow \text{LinRel}_{k((x))}$$

defined as follows



# Interacting Hopf Algebras

The theory  $\mathbb{I}\mathbb{H}$  is  $\text{Circ}$  quotiented by the axioms of  $\mathbb{H}\mathbb{A}$ ,  $\mathbb{H}\mathbb{A}^{op}$  and the following



## Soundness and Completeness

The proof exploits a technique introduced by Steve Lack in *Composing PROPs* (2004)

$$\llbracket c \rrbracket = \llbracket d \rrbracket \iff c \stackrel{\mathbb{I}\mathbb{H}}{=} d$$

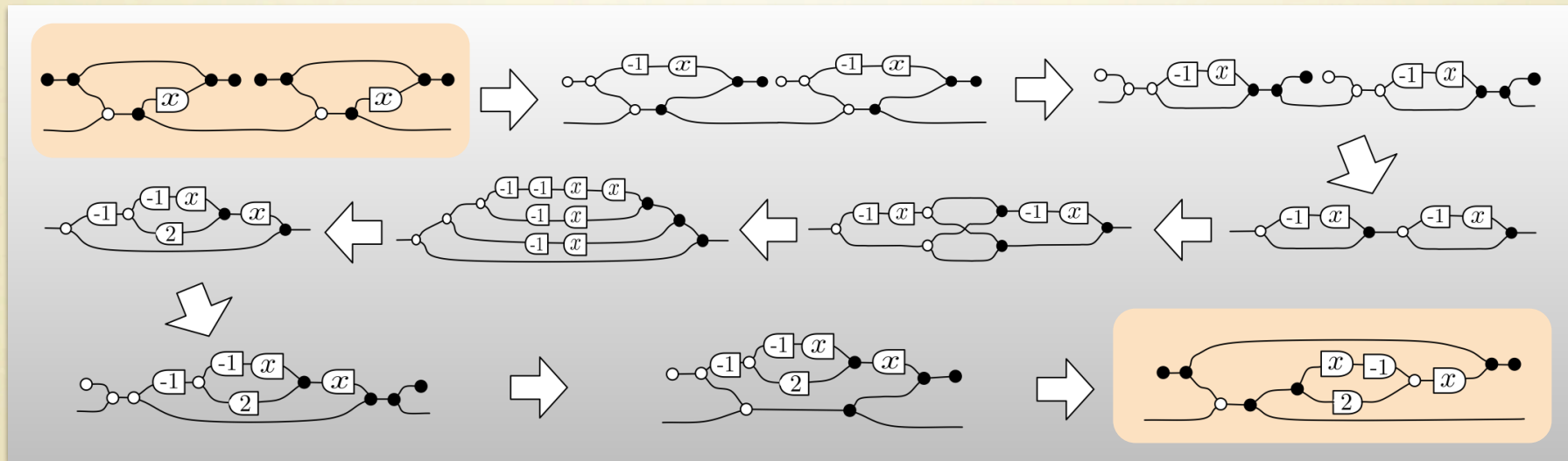
### Kleene's Theorem

$\mathbb{I}\mathbb{H}$  is isomorphic to  $\text{LinRel}_{k(x)}$

Bonchi, Sobocinski, Zanasi  
*Interacting Hopf Algebras*  
 Journal of Pure and Applied Algebra (2017)

# Equational Reasoning

Proof that two diagrams represent the same dynamical system:



Actually, this holds in general:

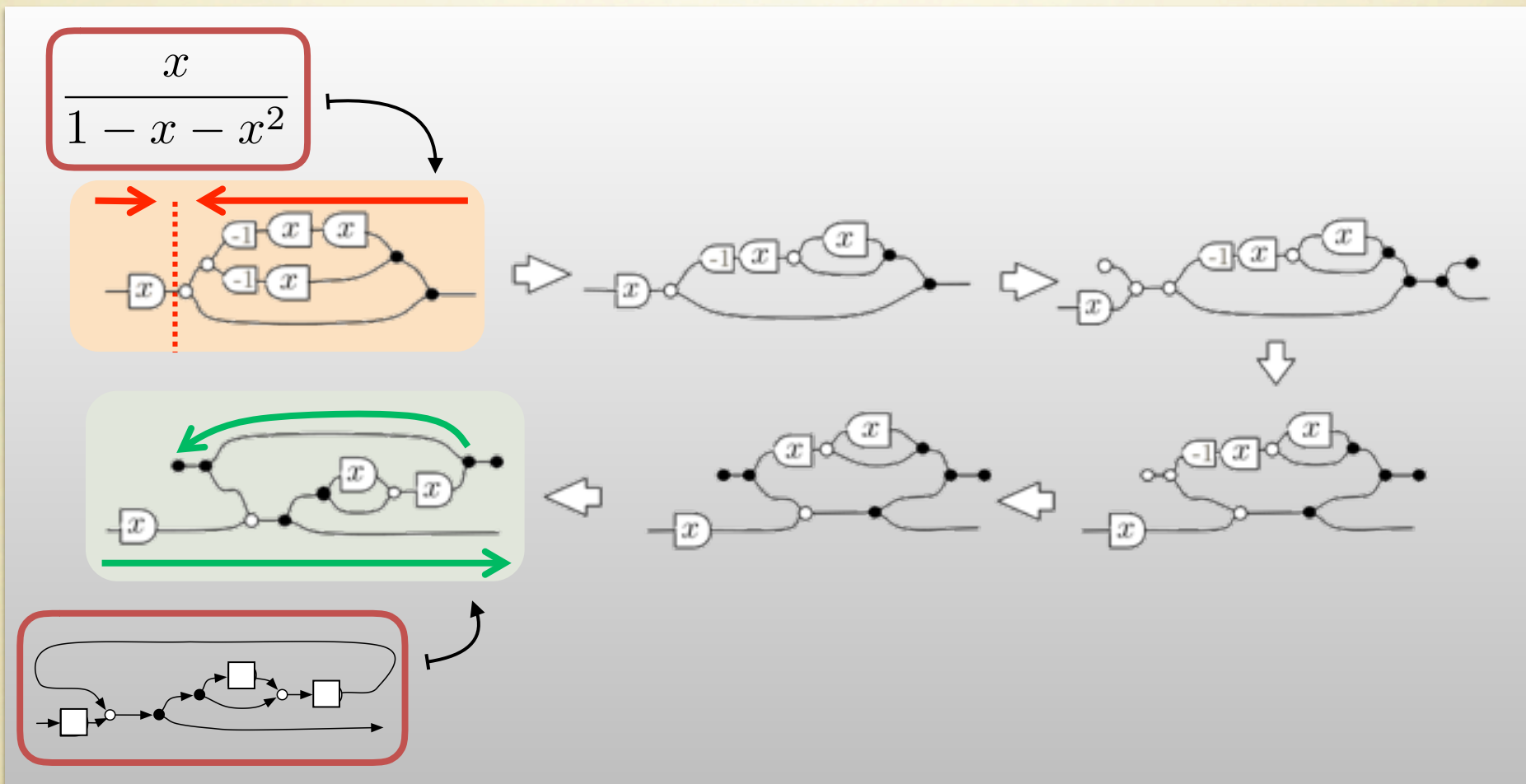
## Normal Form

any circuit is equivalent to one in  $\overrightarrow{\text{Circ}}$  with a feedback loop

The analogue of Bohm-Jacopini's Theorem

# Equational Reasoning

From a specification (a rational fraction) to its implementation as a linear dynamical system.





## **II. CONTROVERSIAL INTERMEZZO**

# Flow direction is FUNDAMENTAL

“flow graphs differ from electrical network graphs in that their branches are directed. In accounting for branch directions it is necessary to take an entirely different line of approach from that adopted in electrical network topology.”

S.J. Mason. *Feedback Theory: I. Some Properties of Signal Flow Graphs*. 1953

# Flow direction is EVIL

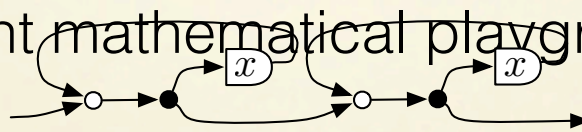
“Adding a signal flow direction is often a figment of one’s imagination, and when something is not real, it will turn out to be cumbersome sooner or later. [...] The input/output framework is totally inappropriate for dealing with all but the most special system interconnections. [The input/output representation] often needlessly complicates matters, mathematically and conceptually. A good theory of systems takes the behavior as the basic notion.”

J. Willems. *Linear systems in discrete time*. 2009

# Flow directionality

Circ does not rely on flow directionality as a primitive

This allows for a more flexible syntax, disclosing a rich and elegant mathematical playground:  $\text{IIIH}$



Whenever flow directionality matters,  
we can always rewrite any circuit in its normal form



*"The reason why physics has ceased to look for causes is that in fact there are no such things. The law of causality, I believe, like much that passes muster among philosophers, is a relic of a bygone age, surviving, like the monarchy, only because it is erroneously supposed to do no harm."*

*(Bertrand Russel - 1913)*

**GRAPHICAL  
LINEAR  
ALGEBRA**

**CARTESIAN AND  
ABELIAN  
BICATEGORIES**

**III. RAMIFICATIONS**

**PETRI  
NETS**

**ELECTRICAL  
CIRCUITS**

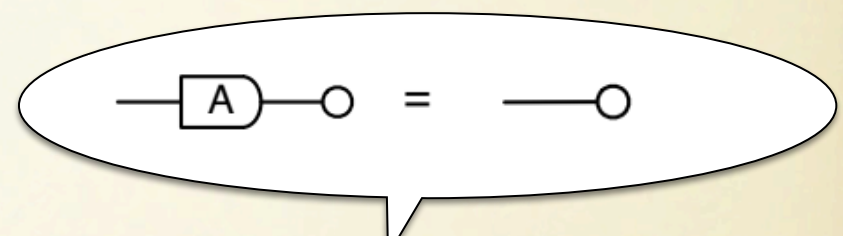
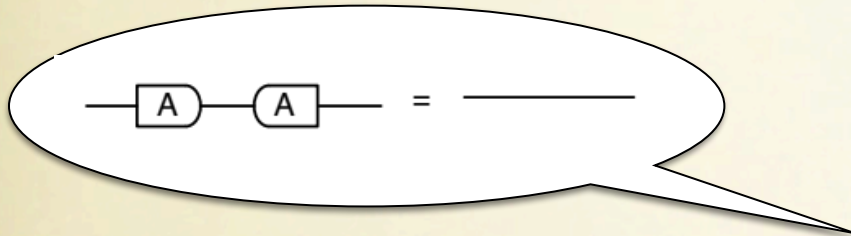
# Graphical Linear Algebra

For dynamical systems, we need the field of fractions of polynomials, but actually the theorem holds for any field

This allows to *reprove*  
well-known theorems of linear algebra  
by means of our axioms

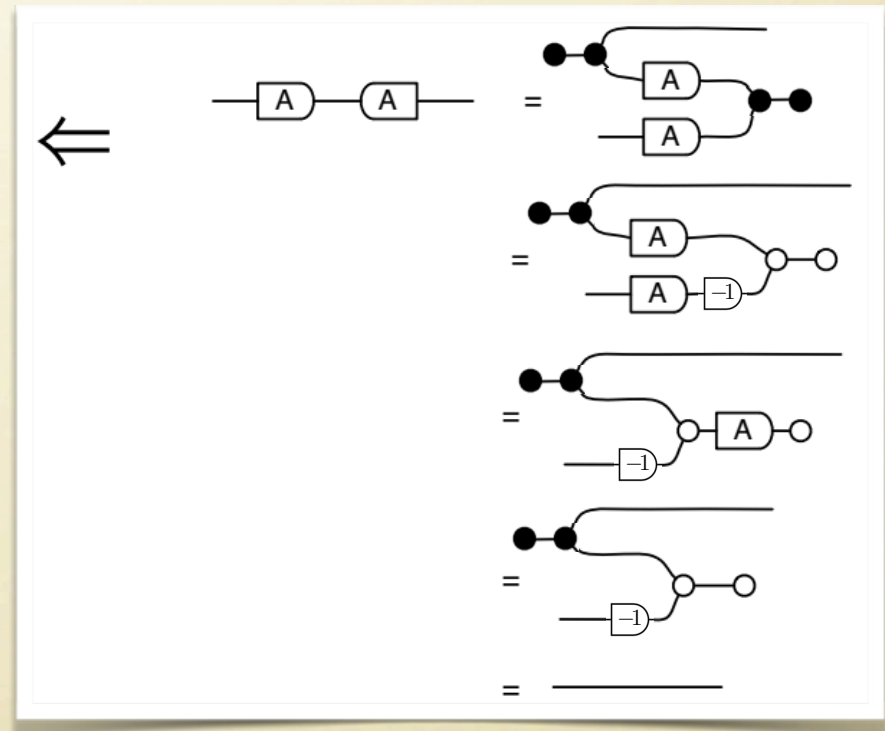
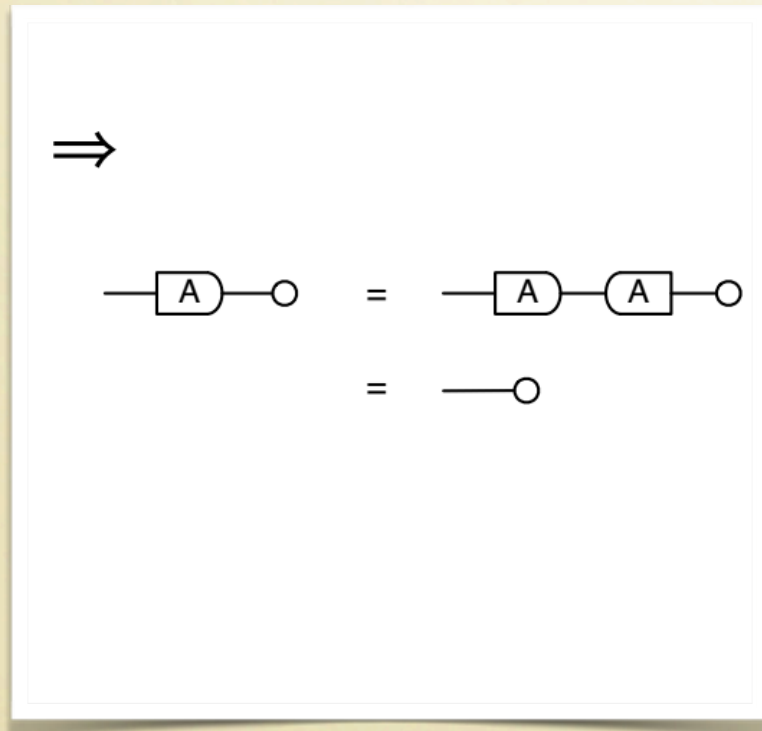
You can find many examples in the blog:  
<https://graphicallinearalgebra.net>

# Equational Proof



**Proposition:** a matrix  $A$  is injective iff its kernel is 0.

## Proof



**GRAPHICAL  
LINEAR  
ALGEBRA**

**CARTESIAN AND  
ABELIAN  
BICATEGORIES**

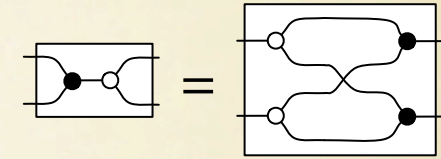
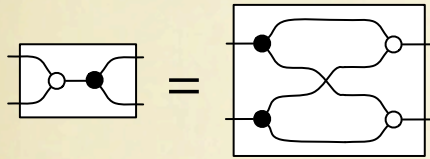
**III. RAMIFICATIONS**

**PETRI  
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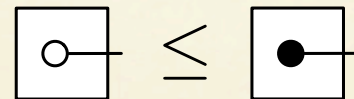


# Cartesian and Abelian Bicategories

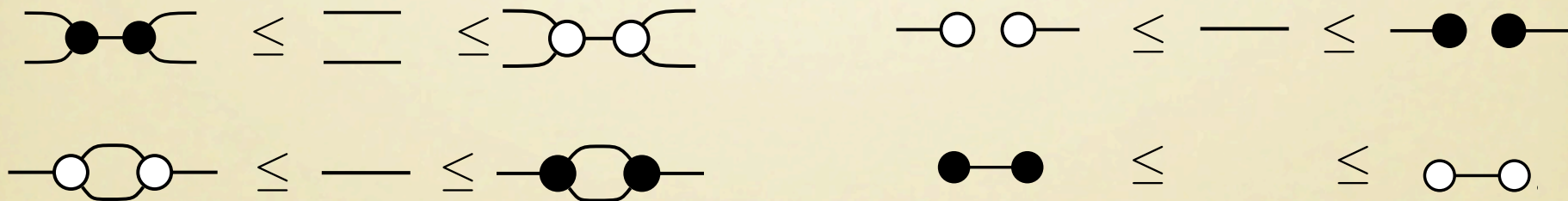


Switching black and white colouring does not change the axiomatisation

To axiomatise the inclusion relation is enough to add



It turns out that  $\mathbb{III}$  is an Abelian bicategory (Carboni-Walters)



**GRAPHICAL  
LINEAR  
ALGEBRA**

**CARTESIAN AND  
ABELIAN  
BICATEGORIES**

**III. RAMIFICATIONS**

**PETRI  
NETS**

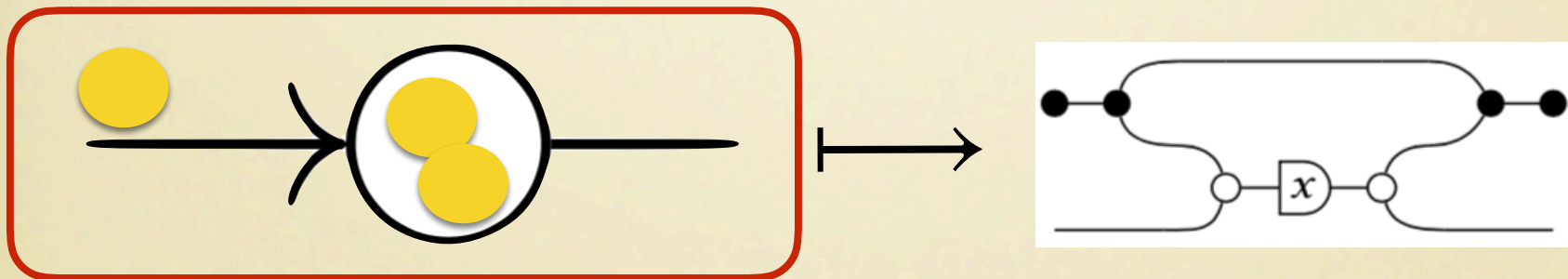
**ELECTRICAL  
CIRCUITS**

# Petri Nets

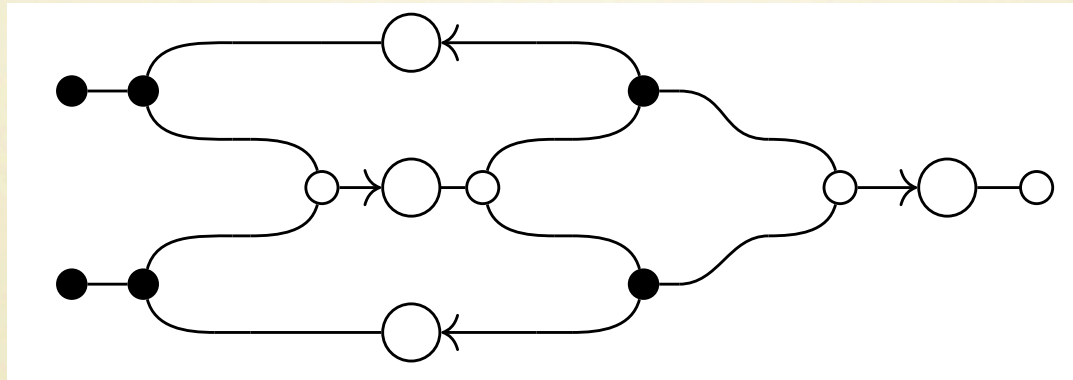
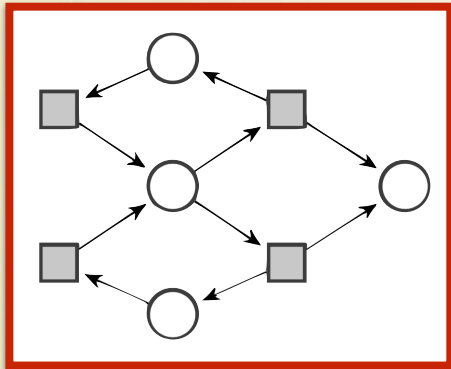
Instead of a field, pick a semiring without additive inverses: the set  $\mathbb{N}$  of natural numbers.

Because  $\mathbb{N}$  does not have additive inverses, it is suitable to model **resources** in a distributed system.

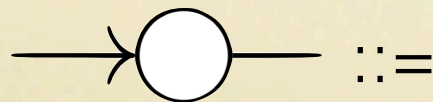
We interpret the **place of a Petri net** as a  $\mathbb{N}$ -circuit diagram



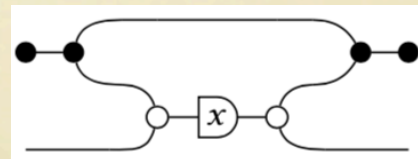
# Petri Nets



dove

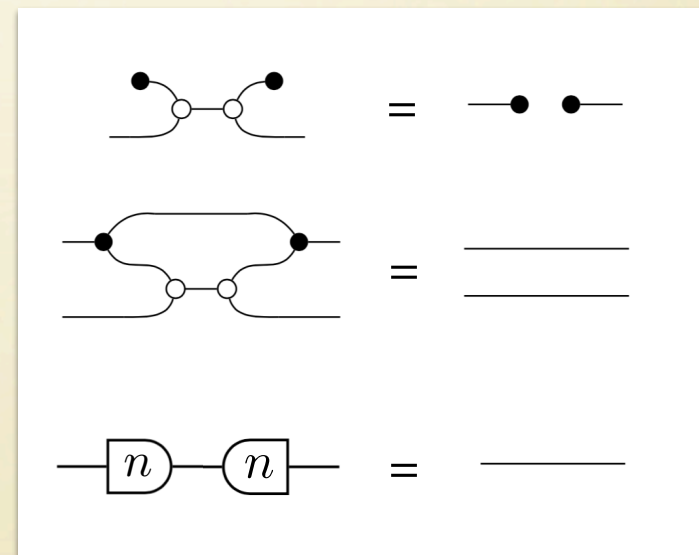
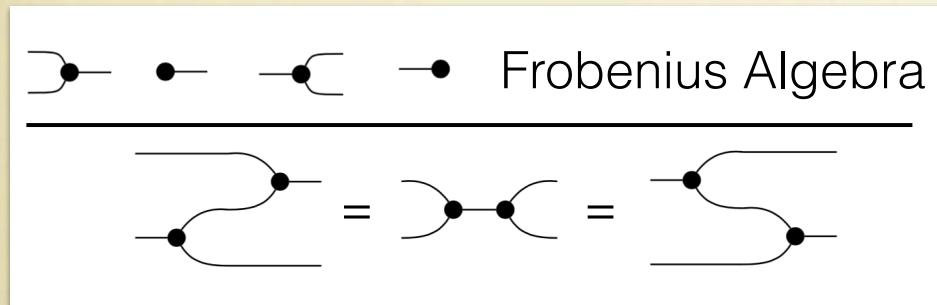
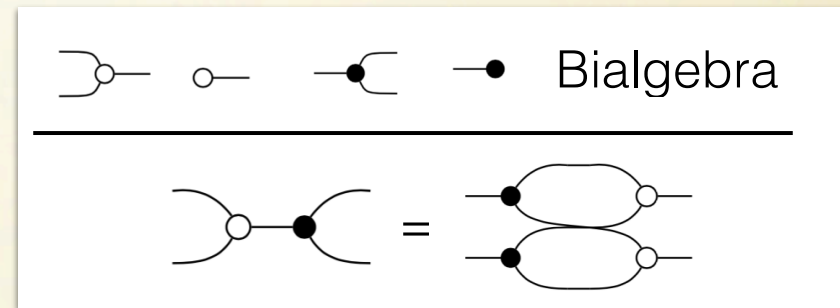
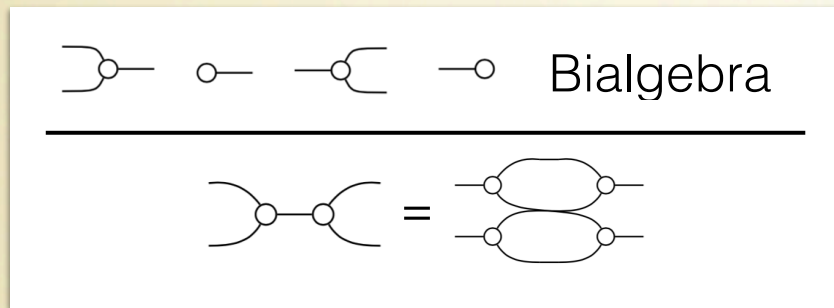


$\equiv$



# Equational Theory

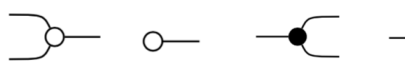
$\mathbb{AR}$ : Algebra of Resources



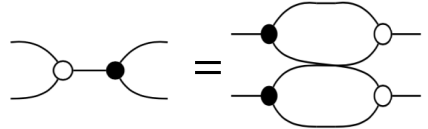

$n \neq 0$

# What are the Fundamental Structures of Concurrency? We still don't know!


Samson Abramsky <sup>1,2</sup>

 Hopf Algebra

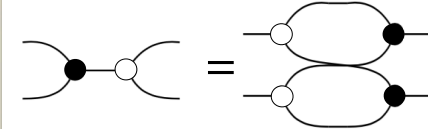

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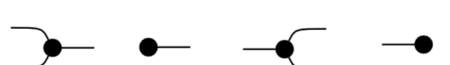
**Not in AR**

 Hopf Algebra

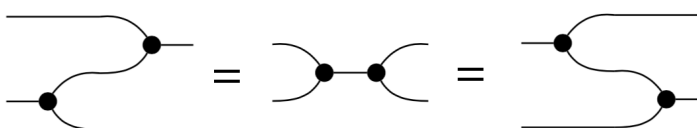
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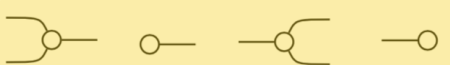
 

**Not in AR**

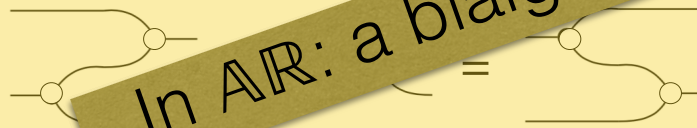
 Frobenius Algebra

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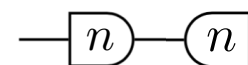



 Frobenius Algebra

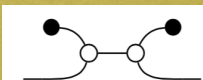

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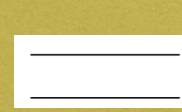


**In AR: a bialgebra**

**Not in AR**

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**GRAPHICAL  
LINEAR  
ALGEBRA**

**CARTESIAN AND  
ABELIAN  
BICATEGORIES**

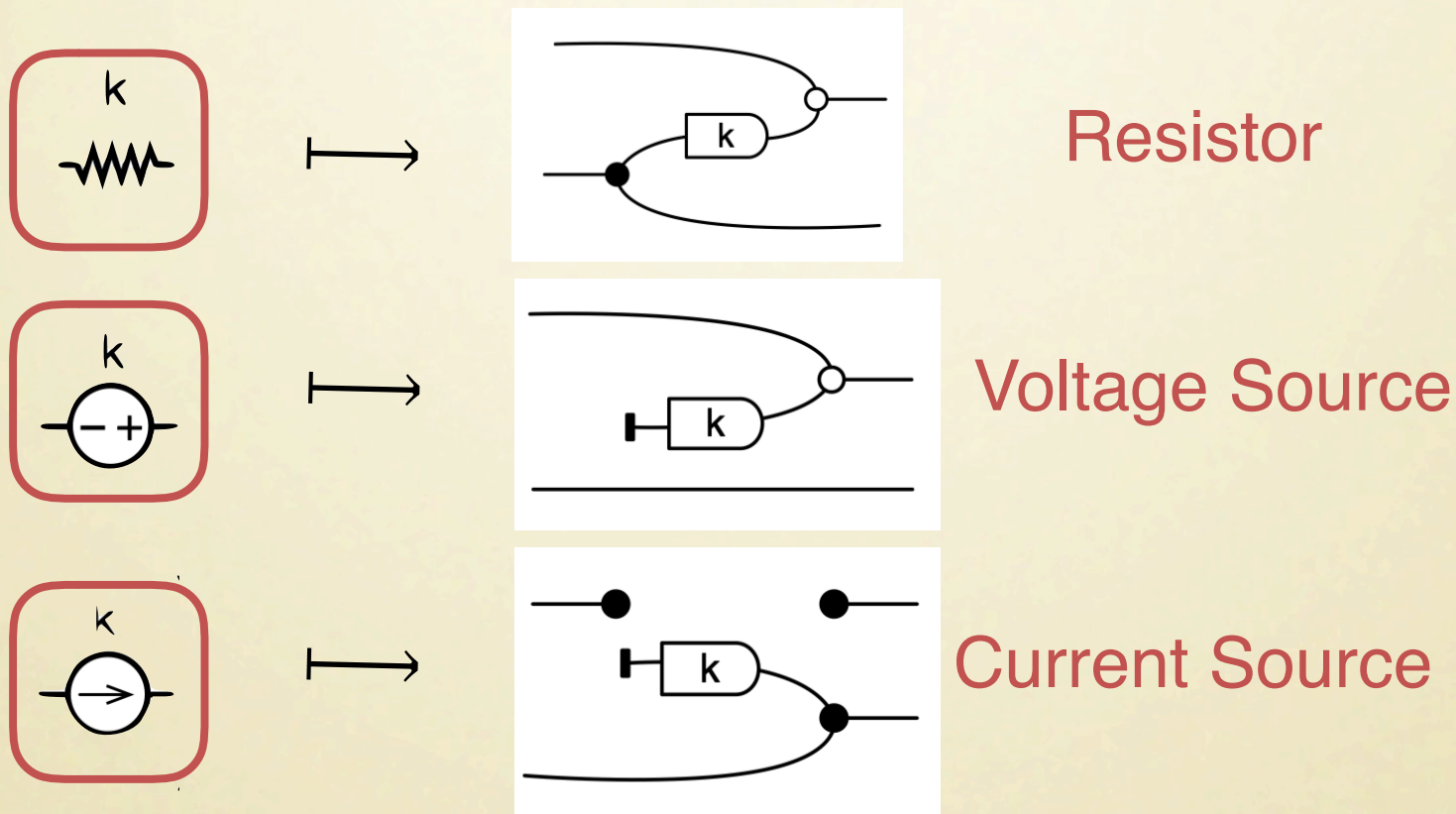
**III. RAMIFICATIONS**

**PETRI  
NETS**

**ELECTRICAL  
CIRCUITS**

# Non Passive Electrical Circuits

Encode electrical circuits as circuit diagrams







# Other References

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- Bonchi, Sobocinski, Zanasi - *The Calculus of Signal Flow Diagrams I: Linear Relations on Streams*, Information and Computation, 2017.

## On Cartesian Bicategories

- Filippo Bonchi, Dusko Pavlovic, Pawel Sobocinski - *Functorial Semantics for Relational Theories*, CoRR abs/1711.08699, 2017.
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